# Discount-based Pricing and Capacity Planning for EV Charging under Stochastic Demand

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# Motivation: Challenges due to widespread EV adoption

- Utilities need to plan capacity robust to load variability
- Our focus: EV Charging Aggregators
  - Buy from utility company, supply to EV users
  - For e.g. parking lots at airports, malls, and DC fast charging facilities etc.
- If an aggregator allows all users to charge at the fastest kW rate, Power capacity  $(K) \propto$  Space  $(M) \times$  Max charging rate  $(r_{\text{max}})$
- For e.g.: UCLA would require  $20,000 \times 120 \text{ kW} = 2.4 \text{ GW}$
- **Problem**: Can we achieve a better scaling than the worst-case product above?
- Solution: Incentivise EV users to stay longer via *pricing*
- $\bullet\,\approx\,10$  fold reduction for aggregator at the scale of UCLA



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$$\underbrace{u_i = \underset{u \ge 0}{\operatorname{argmin}}}_{u \ge 0} \quad \underbrace{P_{a,c,\tau}(x_i, u)}_{\operatorname{Monetary cost}} + \underbrace{\alpha_i \, u}_{\operatorname{Opportunity cost}}$$

• Note: Aggregator could be distributed since sum of Poisson processes is a Poisson process

#### Discount-based pricing function

- Desired properties from  $P(x, \cdot)$  to incentivise longer deadlines  $u_i$ 
  - Decreasing: Users pay lower price for longer service times
  - E.g. Patient users (i.e. lower  $\alpha_i$ , longer  $u_i$ ) get more discount
  - Convex: Discounts are diminishing as service time increases
- We consider the following pricing function
- For charging an EV by x kW-hr in time u hr is,

$$P_{a,c,\tau}(\boldsymbol{x},\boldsymbol{u}) = \boldsymbol{x}(ae^{-\boldsymbol{u}/\tau} + c)$$

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# Discount-based pricing function



Interpreting the parameters:

$$P(\mathbf{x}, \mathbf{u}) = \mathbf{x}(ae^{-\mathbf{u}/\tau} + c)$$

- c Base price (\$/kW-hr),
- a Surge price (\$/kW-hr)
- $\tau$  Suggested service time (hr)

• Service time decision by user *i* to minimize total cost,

$$\underline{u_i} = \tau \cdot \log\left(\frac{ax_i}{\tau \alpha_i}\right)$$

• \$ amount paid by user *i* with demand  $x_i$  and impatience  $\alpha_i$ 

$$P_{a,c,\tau}(x_i, u_i) = c \, x_i + \tau \, \alpha_i$$

#### Capacity planning under stochastic demand

- Let  $f_{x,\alpha}$  be such that  $\mathbb{E}u = b\tau$ ,  $\mathbb{E}\left(\frac{x}{u}\right) = \frac{\mu}{\tau}$  and  $\operatorname{Var}\left(\frac{x}{u}\right) = \frac{1}{2}\left(\frac{\nu}{\tau}\right)^2$
- Let the maximum rate of charging a single EV be  $r_{\text{max}}$  (in kW)

Theorem (Space and Power Capacity, for constant power charging)

For an arrival rate of  $\lambda$ , at any time, with 99% confidence, we have

**1** Demand for space will not exceed

$$M = \lambda \, b\tau + \sqrt{5\lambda \, b\tau} + 3$$

2 Power delivered will not exceed

$$K = \lambda b\mu + (\mu + \nu)\sqrt{6\lambda b/\tau} + 6r_{\max}$$

- Similar to  $Y \sim \mathcal{N}(\mu, \sigma^2)$ , then with 99% confidence,  $Y \leq \mu + 3\sigma$ .
- $K \propto \mathcal{O}(M + r_{\max})$  instead of  $K \propto \mathcal{O}(M \times r_{\max})$
- Optimal rate with respect to  $\lambda$ . A lower bound of  $K = \Omega(\lambda)$  exists.
- **Tradeoff:** Space capacity  $M = \widetilde{\mathcal{O}}(\tau)$ . Power capacity  $K = \widetilde{\mathcal{O}}(1/\sqrt{\tau})$ .

#### Numerical simulations - Excess capacity v/s $\lambda$

Excess capacity  $\propto \frac{1}{\sqrt{\lambda}}$ Excess :=  $\frac{\text{Theorem} - \text{Actual}}{\text{Actual}} \times 100$ . For 100 simulations of 8 hr each.



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#### Numerical simulations - Capacity v/s Confidence



Figure:  $N^{\delta}$  is the 1 -  $\delta$  percentile of occupied space,  $Q^{\delta}$  is the 1 -  $\delta$  percentile of power delivery.



- Pricing can help incite some desired behaviour in *impatient* users
- Probabilistic constraints allow reducing installation capacity drastically from  $K \propto \mathcal{O}(M \times r_{\text{max}})$  to  $K \propto \mathcal{O}(M + r_{\text{max}})$ 
  - Example:  $\approx 10$  fold reduction for UCLA.
- We characterized probability of failure that helps plan for back-up capacity (battery banks/generators)
- Other aggregators with impatient users
  - *Cloud computing*: Client machines upload FLOPS to a server with a deadline for the computation. Users get discounts for waiting longer.
  - *Cab aggregators*: Users get a discount for waiting longer (Uber Pool or Lyft Line), allowing for more efficient resource allocation by aggregator.

# Thank You!